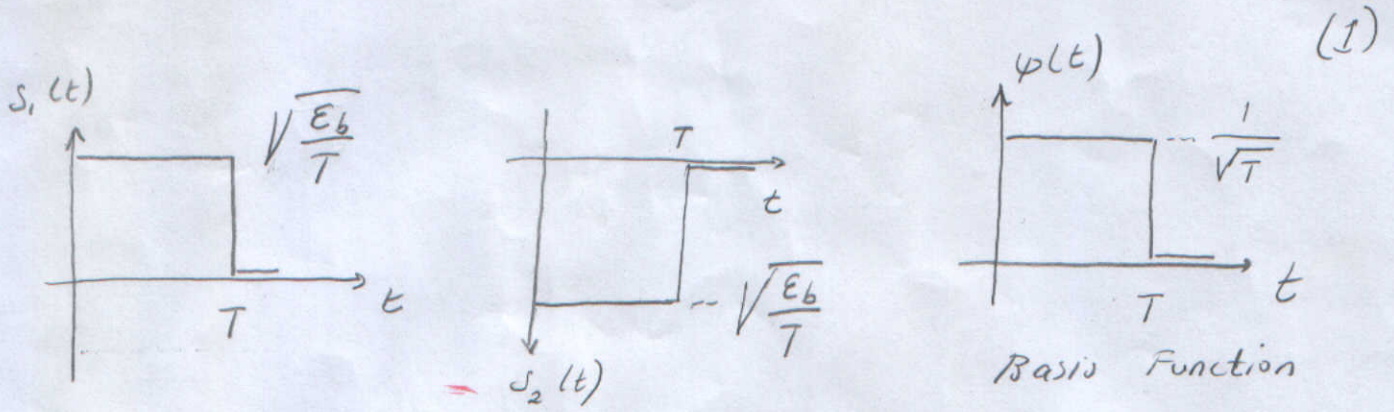


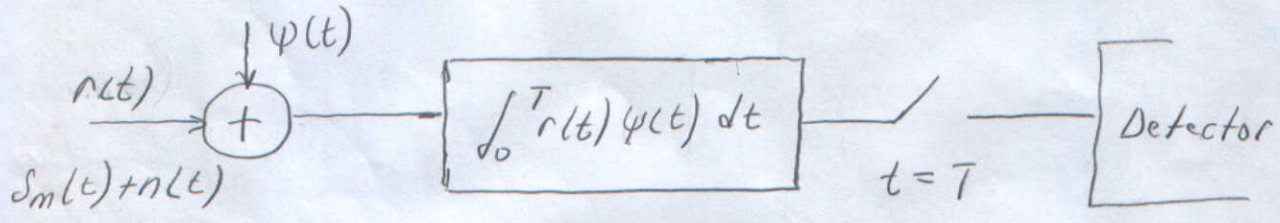
Formulating Probability of Error for Binary PAM



$$s_1(t) = \sqrt{\frac{E_b}{T}} \quad 0 < t < T, \quad s_2(t) = -\sqrt{\frac{E_b}{T}} \quad 0 < t < T$$

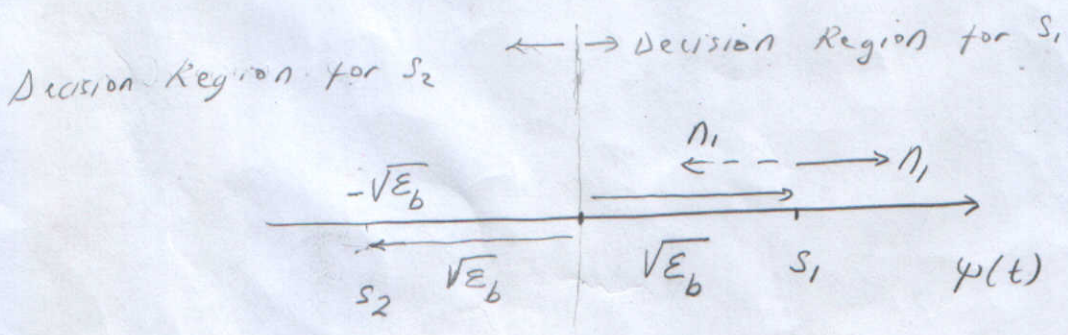
$$s_1 = [\sqrt{E_b}] \quad s_2 = [-\sqrt{E_b}] \quad , \quad E_b = \int_0^T s_m^2(t) dt$$

Block diagram of Matched Filter Correlator



Assume s_1 was transmitted

$$\int_0^T r(t)\phi(t) dt = \sqrt{E_b} + n_1 = r \text{ (vector)}$$



(2)

Since it is one dimensional space n_1 will always be along $\psi(t)$ axis. Assuming that $s_1(t)$ and $s_2(t)$ have the representation above having equal energy.

Thus the threshold of decision is set to zero

Then error will result if $n_1 \leq -\epsilon_b^{1/2}$

The same result could be obtained via $C(r, s_m)$ metric. In the above case of s_1 being transmitted

$$\begin{aligned} C(r, s_1) &= 2r \cdot s_m - \|s_m\|^2 \\ &= 2[\sqrt{\epsilon_b} + n_1] \cdot [\epsilon_b^{1/2}] - \epsilon_b \\ &= \epsilon_b + 2n_1 \epsilon_b^{1/2} \end{aligned}$$

$$\begin{aligned} C(r, s_2) &= 2[\epsilon_b^{1/2} + n_1] [-\epsilon_b^{1/2}] - \epsilon_b \\ &= -3\epsilon_b - 2n_1 \epsilon_b^{1/2} \end{aligned}$$

For error to occur $C(r, s_2) > C(r, s_1)$ or

$$-3\epsilon_b - 2n_1 \epsilon_b^{1/2} > \epsilon_b + 2n_1 \epsilon_b^{1/2}$$

$$n_1 < -\epsilon_b^{1/2}$$

Note that due to symmetry of Gaussian distribution

The definition of P_e is based on Eq (7.5.55)

that is

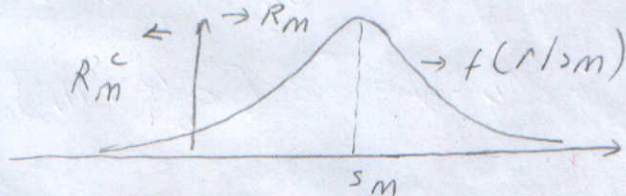
$$P(e/s_m) = \int_{R_m^c} f(r/s_m) dr$$

R_m means the region of correct decision

R_m^c means the region outside the region of

correct decision, i.e. it is the complementary to

R_c



Note that

$$f(r/s_m) = \frac{1}{(\pi N_0)^{1/2}} e^{-\frac{(r-s_m)^2}{N_0}}, \text{ hence}$$

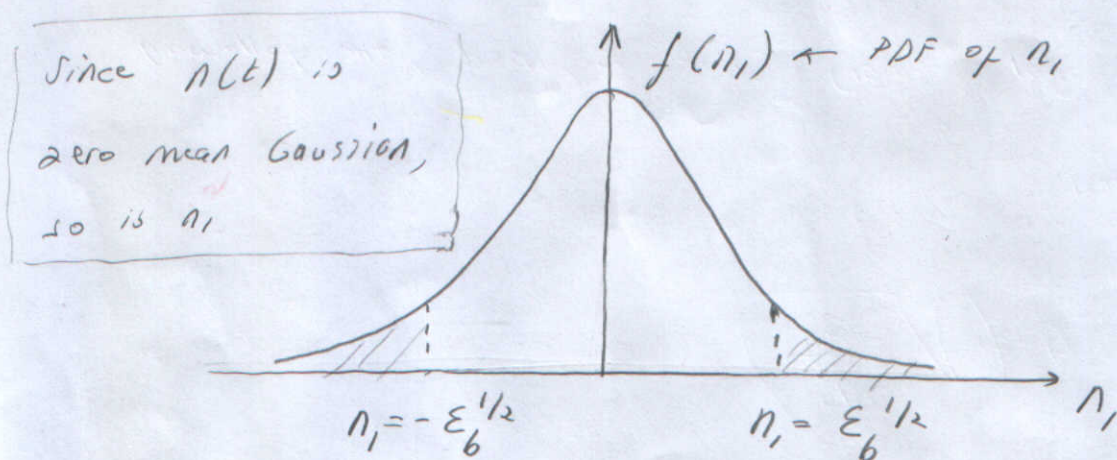
$$\int_{-\infty}^0 \frac{1}{(\pi N_0)^{1/2}} e^{-\frac{(r-s_m)^2}{N_0}} dr = \frac{1}{(\pi N_0)^{1/2}} \int_{-\infty}^{-\sqrt{E_b}} e^{-\frac{x^2}{N_0}} dx$$

Same as n_1 integral

(3)

$$-\infty < n_1 < -\epsilon_b^{1/2} \quad \text{and} \quad \epsilon_b^{1/2} < n_1 < \infty$$

will give the same area (see the figure below)



$$f(n_1) = \frac{1}{(\pi N_0)^{1/2}} e^{-n_1^2/N_0}$$

$$P_e = \text{Prob}(\epsilon_b^{1/2} < n_1 < \infty) \leftarrow \text{(This is } P(e/s_2))$$

$$= \int_{\sqrt{\epsilon_b}}^{\infty} \frac{1}{(\pi N_0)^{1/2}} e^{-n_1^2/N_0} dn_1$$

Make the substitution $\frac{n_1^2}{N_0} = \frac{z^2}{2}$

$$P_e(e/s, \text{ binary PAM}) = \frac{1}{(2\pi)^{1/2}} \int_{\sqrt{\frac{2\epsilon_b}{N_0}}}^{\infty} e^{-z^2/2} dz$$

$$P_e = Q\left(\sqrt{\frac{2\epsilon_b}{N_0}}\right) \text{ same as Eq (7.6.8) of Proakis}$$

The above is the same as binary PSK with antipodal signalling

If orthogonal binary PSK is used, the following is obtained

$$P_e = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

hence P_e (orthogonal binary PSK)

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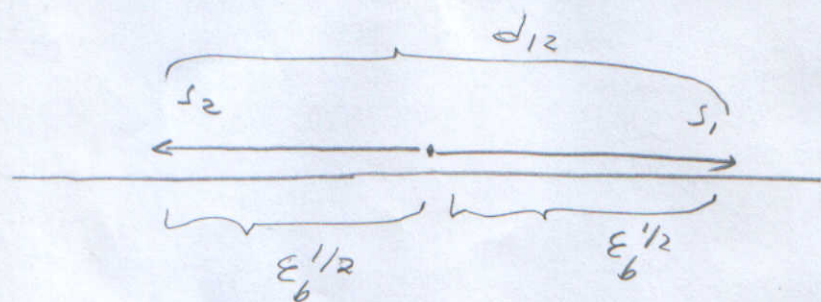
P_e (antipodal binary PSK)

4.05.2005

Addendum for orthogonal binary PSK (1a)

Note that the above $P_e = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$ is

also valid for binary PSK where $s_1(t)$ and $s_2(t)$ are antipodal signals, i.e.,



Of course $s_1(t)$ and $s_2(t)$ must be equiprobable.

We observe that in Q. function E_b may also be expressed in terms of the distance between end of vectors s_1 and s_2 . In this manner

$$d_{12} = \sqrt{E_b} + \sqrt{E_b} = 2\sqrt{E_b}$$

$$2E_b = d_{12}^2/2 \quad \text{or} \quad d_{12}^2 = 4E_b$$

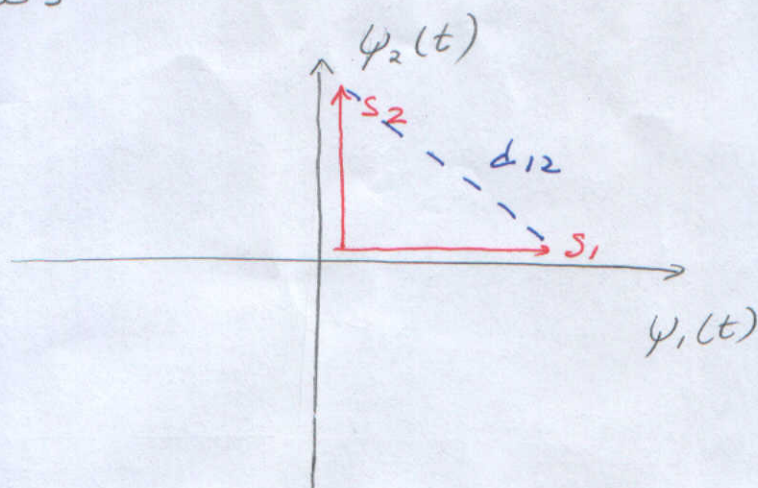
Hence P_e may be written as

$$P_e = Q\left(\sqrt{\frac{d_{12}^2}{2N_0}}\right)^{2E_b}$$

See also Proakis 2000 pp. 406-407

(2a)

Therefore by taking the binary orthogonal PSK as follows



$$d_{12} \text{ (for orthogonal binary PSK)} = \sqrt{2E_b}$$

or $d_{12}^2 = 2E_b$, thus

$$P_e \text{ (for orthogonal binary PSK)} = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

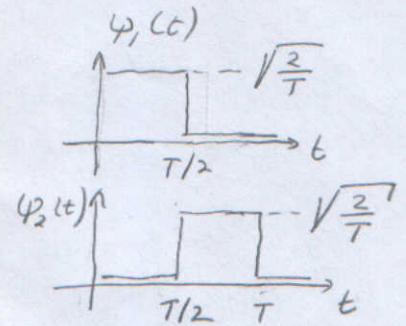
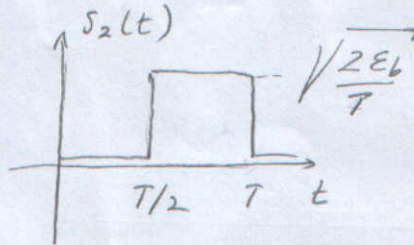
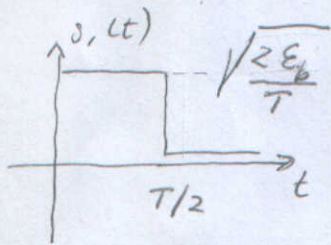
As expected P_e of orthogonal case is has larger values than P_e of antipodal case.

The above result may have been achieved

by examining the correlation metrics values.

(3a)

For this case, the waveform illustrations are reproduced below



$$\text{So } s_1(t) = \sqrt{E_b} \psi_1(t) + 0 \psi_2(t) \quad s_1 = [\sqrt{E_b}; 0]$$

$$s_2(t) = 0 \psi_1(t) + \sqrt{E_b} \psi_2(t) \quad s_2 = [0; \sqrt{E_b}]$$

This way the OP from the correlators (or VAF)

If s_1 was transmitted

$$r_1 = \sqrt{E_b} + n_1, \quad r_2 = n_2 \quad \text{or } r = [\sqrt{E_b} + n_1, n_2]$$

Normally test against correlator metrics for s_1 and s_2 should yield that $C(r, s_1) > C(r, s_2)$

So in case of error the opposite should

occur, that is $C(r, s_2) > C(r, s_1)$

This is shown below

(4a)

$$\begin{aligned}
 C(r, s_1) &= 2r \cdot s_1 - \|s_1\|^2 \\
 &= 2[\sqrt{\epsilon_b} + n_1, n_2] \begin{bmatrix} \sqrt{\epsilon_b} \\ 0 \end{bmatrix} - \epsilon_b \\
 &= 2(\epsilon_b + \epsilon_b^{1/2} n_1) - \epsilon_b
 \end{aligned}$$

$$\begin{aligned}
 C(r, s_2) &= 2r \cdot s_2 - \|s_2\|^2 \\
 &= 2[\sqrt{\epsilon_b} + n_1, n_2] \begin{bmatrix} 0 \\ \sqrt{\epsilon_b} \end{bmatrix} - \epsilon_b \\
 &= 2n_2 \epsilon_b^{1/2} - \epsilon_b
 \end{aligned}$$

The condition of $C(r, s_2) \geq C(r, s_1)$

$$\sqrt{\epsilon_b} (2n_2 - 1) \geq \sqrt{\epsilon_b} (2\sqrt{\epsilon_b} + 2n_1 - 1)$$

$$2n_2 - 2n_1 \geq 2\sqrt{\epsilon_b}$$

$$n_2 - n_1 \geq \epsilon_b^{1/2}$$

n_1 and n_2 are independent Gaussian variables

with zero mean and each with $N_0/2$, hence

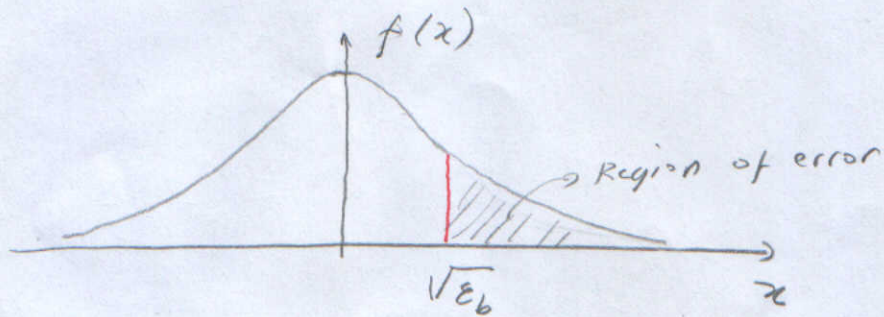
$x = n_2 - n_1$ is again with zero mean, but with

No variance ($N_0/2 + N_0/2 = N_0$)

(5a)

Hence

$$f(n_2 - n_1 = x) = \frac{1}{\sqrt{2\pi N_0}} e^{-x^2/2N_0}$$



P_e (of orthogonal binary PSK) = $f(x)$ being larger than \sqrt{E}

$$= \frac{1}{\sqrt{2\pi N_0}} \int_{\sqrt{E_b}}^{\infty} e^{-x^2/2N_0}$$

$$= \frac{1}{\sqrt{2\pi}} \int_{\sqrt{E_b/N_0}}^{\infty} e^{-x^2} dx$$

$$= Q \left[\sqrt{\frac{E_b}{N_0}} \right]$$

The same result as before